

TD 2: Polynomial multiplication

Exercise 1. Polynomial Arithmetic

A polynomial $P = p_0 + p_1X + \dots + p_{n-1}X^{n-1}$ by an array of n elements $[p_0, p_1, \dots, p_{n-1}]$.

a. Propose a *Divide and Conquer* algorithm for the multiplication of two polynomials of same degree.

Indication: one can use the identity:

$$(P_0 + XP_1)(Q_0 + XQ_1) = P_0Q_0 + X(P_0Q_1 + P_1Q_0) + X^2P_1Q_1$$

b. Analyse its cost (when the size of the polynomials is a power of two).

c. In 1960, Karatsuba proposed to use instead the following formula:

$$(P_0 + XP_1)(Q_0 + XQ_1) = P_0Q_0 + X((P_0 + P_1)(Q_0 + Q_1) - P_0Q_0 - P_1Q_1) + X^2P_1Q_1$$

Deduce an algorithm and analyse its complexity.

We now investigate the decomposition of polynomials in three: $P = P_0 + XP_1 + X^2P_2$ et $Q = Q_0 + XQ_1 + X^2Q_2$. Toom proposed a formula computing $P \times Q$ using the five following values:

$$\begin{aligned} M_0 &= P_0Q_0 \\ M_1 &= (P_0 + P_1 + P_2)(Q_0 + Q_1 + Q_2) \\ M_2 &= (P_0 - P_1 + P_2)(Q_0 - Q_1 + Q_2) \\ M_3 &= (P_0 + 2P_1 + 4P_2)(Q_0 + 2Q_1 + 4Q_2) \\ M_4 &= P_2Q_2 \end{aligned}$$

The product $R = P \times Q = R_0 + R_1X + R_2X^2 + R_3X^3 + R_4X^4$ is obtained as follows:

$$\begin{cases} R_0 = M_0 \\ R_1 = \frac{1}{6}(-3M_0 + 6M_1 - 2M_2 - M_3 + 12M_4) \\ R_2 = \frac{1}{2}(-2M_0 + M_1 + M_2 - 2M_4) \\ R_3 = \frac{1}{6}(3M_0 - 3M_1 - M_2 + M_3 - 12M_4) \\ R_4 = M_4 \end{cases} \quad (1)$$

d. What is the cost of the corresponding algorithm multiplying polynomials of arbitrary degrees?

e. Justify that the formulas computing the M_i can be viewed as evaluations. State in which points?

f. Deduce how the coefficients of the formule (1) have been found.

More generally, Toom-Cook algorithms at order k compute the product $P \times Q$ of two polynomials of size k in $(2k - 1)$ multiplications.

g. What is the cost of these algorithms ?

h. Explain how to construct them.

i. Conclude on the cost of multiplying polynomials.